

Games, graphs, and machines

In **tropical algebra**, the sum of two numbers is their minimum and the product of two numbers is their sum.

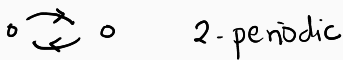
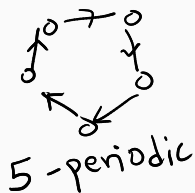
August 19, 2025

Periodic?

Let A be the boolean adjacency matrix of a graph.

Claim — The sequence of matrices A, A^{*2}, A^{*3}, \dots is periodic.

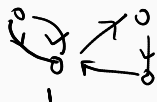
True or false?



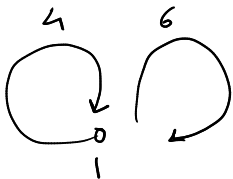
} Ex



$$A^1 = \begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix} \quad A^2 = 0 \quad A^3 = 0 \dots$$



A^{*n}	n	1	2	3	4	5	6	7	...
$(1,1)$		0	1	1	1	1	1	1	...



n	1	2	3	4	5	6	7	8	9	10	11	12	13	14
(1,1)	0	0	0	1	0	1	0	1	0	1	0	1	0	1

"eventually periodic"

Thm.: $A^0, A^{*1}, A^{*2}, A^{*3}, \dots, A^{*7}, A^{*8}$ is eventually periodic.

Proof: $n \times n$ boolean matrices \leftarrow finite (2^{n^2})

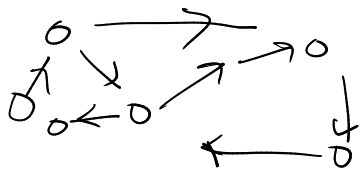
\Rightarrow There is a repeat (∞ seq.)

$$A^{*n_1} = A^{*(n_1+N)}$$

Chaos



Given a graph \rightsquigarrow period of A^n
(number)



↓
99
—

Warm up with min/plus arithmetic

Remember that $\oplus = \min$ and $\odot = +$. Find

Numbers

"

$\mathbb{R} \cup \{+\infty\}$

\oplus, \odot obey the laws of $+, \cdot$.

$$3 : \underbrace{2 \odot (3 \oplus 1)}_3 \oplus \underbrace{3 \odot (\infty \oplus 2)}_5$$

$$a \odot (b \oplus c) = a \odot b \oplus a \odot c$$

Boolean

$+$: OR

\cdot : AND

Tropical

\oplus cost of OR.

\odot cost of AND

Tropical

\oplus

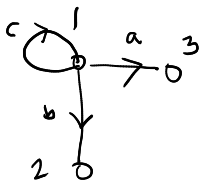
identity :

∞

\odot

identity :

0



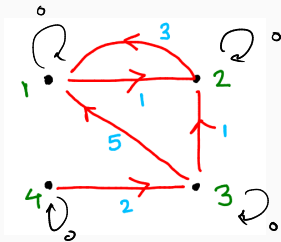
$$\begin{pmatrix} c & b & a \\ \infty & & \end{pmatrix}$$

adj matrix,
"cost
matrix"

Weighted adjacency matrix

Write the min/plus weighted adjacency matrix of the graph. Assume that the loops have weight 0 (not shown).

$$A = \begin{pmatrix} 0 & 1 & \infty & \infty \\ 3 & 0 & \infty & \infty \\ 5 & 1 & 0 & \infty \\ \infty & \infty & 2 & 0 \end{pmatrix}$$



$$\begin{matrix} 0 & 1 & \infty & \infty \\ 3 & 0 & \infty & \infty \\ 4 & 1 & 0 & \infty \\ 7 & 3 & 2 & 0 \end{matrix}$$

Min/plus powers

Find the min/plus square and ~~cube~~ of the adjacency matrix. What do its entries represent?

$$\begin{pmatrix} 0 & 1 & \infty & \infty \\ 3 & 0 & \infty & \infty \\ 5 & 1 & 0 & \infty \\ \infty & \infty & 2 & 0 \end{pmatrix} \begin{pmatrix} 0 & 1 & \infty & \infty \\ 3 & 0 & \infty & \infty \\ 5 & 1 & 0 & \infty \\ \infty & \infty & 2 & 0 \end{pmatrix} = \begin{pmatrix} 0 & 1 & \infty & \infty \\ 3 & 0 & \infty & \infty \\ 4 & 1 & 0 & \infty \\ 7 & 3 & 2 & 0 \end{pmatrix}$$

$A^{\odot 2}$ = Shortest (least cost) paths of length 2.

$A^{\odot 3}$ = Least costly paths of length 3

$A^{\odot 3}, A^{\odot 4}, A^{\odot 5}, \dots$ stabilises.

Why do min/plus powers give shortest paths?

For example, the third power:

$$\begin{aligned}A_{i,j}^3 &= (A_{i,1}^2 \odot A_{1,j}) \oplus (A_{i,2}^2 \odot A_{2,j}) \oplus (A_{i,3}^3 \odot A_{3,j}) \oplus (A_{i,4}^2 \odot A_{4,j}) \\ &= \min(A_{i,1}^2 + A_{1,j}, A_{i,2}^2 + A_{2,j}, A_{i,3}^3 + A_{3,j}, A_{i,4}^2 + A_{4,j})\end{aligned}$$

When do we stop?

Assume:

1. we have all loops with weight 0,
2. all weights are non-negative.

Theorem

Let n be the number of vertices. Then $A^{\odot(n-1)} = A^{\odot n} = A^{\odot(n+1)} = \dots$.